

# On some self-orthogonal codes from $M_{11}$

Ivona Novak

(inovak@math.uniri.hr)

joint work with Vedrana Mikulić Crnković

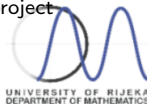
(vmikulic@math.uniri.hr)

Department of Mathematics, University of Rijeka

Finite Geometry & Friends,

A Brussels Summer School on Finite Geometry




This work has been supported by Croatian Science Foundation under the project 6732 and by the University of Rijeka under the project number uniri-prirod-18-111-1249.



Weakly self-orthogonal designs from  $M_{11}$

Codes from  $M_{11}$

Codes from orbit matrices of weakly  $q$ -self-orthogonal 1-designs

-  V. Tonchev, Self-Orthogonal Designs and Extremal Doubly-Even Codes, *Journal of Combinatorial Theory, Series A* 52, 197-205 (1989).
-  D. Crnković, V. Mikulić Crnković, A. Svob, On some transitive combinatorial structures constructed from the unitary group  $U(3, 3)$ , *J. Statist. Plann. Inference* 144 (2014), 19-40.
-  D. Crnković, V. Mikulić Crnković, B.G. Rodrigues, On self-orthogonal designs and codes related to Held's simple group, *Advances in Mathematics of Communications* 607-628 (2018).

## Mathieu group $M_{11}$

$M_{11}$  is simple group of order 7920 which has 39 non-equivalent transitive permutation representations.

Among others, lattice of  $M_{11}$  is consisted of 1 subgroup of index 22, 1 subgroup of index 55, 1 subgroup of index 66, 3 subgroups of index 110, 2 subgroups of index 132, 1 subgroup of index 144 and 1 subgroup of index 165. Subgroup of  $M_{11}$  with largest index has index 3960.

Using mentioned subgroups we obtained transitive permutation representations of  $M_{11}$  on 22, 55, 66, 110, 132, 144 and 165 points.

## Weakly self-orthogonal designs

An incidence structure  $\mathcal{D} = (\mathcal{P}, \mathcal{B}, \mathcal{I})$ , with point set  $\mathcal{P}$ , block set  $\mathcal{B}$  and incidence  $\mathcal{I}$  is called a  $t - (v, k, \lambda)$  design, if  $\mathcal{P}$  contains  $v$  points, every block  $B \in \mathcal{B}$  is incident with  $k$  points, and every  $t$  distinct points are incident with  $\lambda$  blocks.

The incidence matrix of a design is a  $b \times v$  matrix  $[m_{ij}]$  where  $b$  and  $v$  are the numbers of blocks and points respectively, such that  $m_{ij} = 1$  if the point  $P_j$  and the block  $B_i$  are incident, and  $m_{ij} = 0$  otherwise.

A design is weakly  $q$ -self-orthogonal if all the block intersection numbers gives the same residue modulo  $q$ .

A weakly  $q$ -self-orthogonal design is  $q$ -self-orthogonal if the block intersection numbers and the block sizes are multiples of  $q$ .

Specially, weakly 2-self-orthogonal design is called weakly self-orthogonal design, and 2-self-orthogonal design is called self-orthogonal.

## Construction

### Theorem ([2])

Let  $G$  be a finite permutation group acting transitively on the sets  $\Omega_1$  and  $\Omega_2$  of size  $m$  and  $n$ , respectively. Let  $\alpha \in \Omega_1$  and  $\Delta_2 = \bigcup_{i=1}^s \delta_i G_\alpha$ , where  $\delta_1, \dots, \delta_s \in \Omega_2$  are representatives of distinct  $G_\alpha$ -orbits. If  $\Delta_2 \neq \Omega_2$  and

$$\mathcal{B} = \{\Delta_2 g \mid g \in G\},$$

then  $\mathcal{D} = (\Omega_2, \mathcal{B})$  is  $1 - (n, |\Delta_2|, \frac{|G_\alpha|}{|G_{\Delta_2}|} \sum_{i=1}^n |\alpha G_{\delta_i}|)$  design with  $\frac{m \cdot |G_\alpha|}{|G_{\Delta_2}|}$  blocks.

Using mentioned construction for transitive permutation representations of  $M_{11}$ , we constructed 169 non-isomorphic weakly self-orthogonal designs:

- ▶ 6 designs on 66 points,
- ▶ 41 designs on 110 points,
- ▶ 76 designs on 132 points,
- ▶ 26 designs on 144 points,
- ▶ 20 designs on 165 points.

Two of constructed designs are 2-designs:  $2 - (144, 66, 30)$  and its complement.

## Codes from weakly self-orthogonal designs

### Theorem ([1])

Let  $\mathcal{D}$  be weakly self-orthogonal design and let  $M$  be it's  $b \times v$  incidence matrix.

- ▶ If  $\mathcal{D}$  is a self-orthogonal design, then the matrix  $M$  generates a binary self-orthogonal code.
- ▶ If  $\mathcal{D}$  is such that  $k$  is even and the block intersection numbers are odd, then the matrix  $[I_b, M, \mathbf{1}]$  generates a binary self-orthogonal code.
- ▶ If  $\mathcal{D}$  is such that  $k$  is odd and the block intersection numbers are even, then the matrix  $[I_b, M]$  generates a binary self-orthogonal code.
- ▶ If  $\mathcal{D}$  is such that  $k$  is odd and the block intersection numbers are odd, then the matrix  $[M, \mathbf{1}]$  generates a binary self-orthogonal code.

Codes from weakly  $q$ -self-orthogonal designs**Theorem**

Let  $q$  be prime power and  $\mathbb{F}_q$  a finite field of order  $q$ . Let  $\mathcal{D}$  be a weakly  $q$ -self-orthogonal design such that  $k \equiv a \pmod{q}$  and  $|B_i \cap B_j| \equiv d \pmod{q}$ , for all  $i, j \in \{1, \dots, b\}$ ,  $i \neq j$ , where  $B_i$  and  $B_j$  are two blocks of a design  $\mathcal{D}$ . Let  $M$  be it's  $b \times v$  incidence matrix.

- ▶ If  $\mathcal{D}$  is  $q$  self-orthogonal design, then  $M$  generates a self-orthogonal code over  $\mathbb{F}_q$ .
- ▶ If  $a = 0$  and  $d \neq 0$ , then the matrix  $[\sqrt{d} \cdot I_b, M, \sqrt{-d} \cdot \mathbf{1}]$  generates a self-orthogonal code over  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-d$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
- ▶ If  $a \neq 0$  and  $d = 0$ , then the matrix  $[M, \sqrt{-a} \cdot I_b]$  generates a self-orthogonal code over  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-a$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
- ▶ If  $a \neq 0$  and  $d \neq 0$ , there are two cases:
  1. if  $a = d$ , then the matrix  $[M, \sqrt{-d} \cdot \mathbf{1}]$  generates a self-orthogonal code over  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-a$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise, and
  2. if  $a \neq d$ , then the matrix  $[\sqrt{d-a} \cdot I_b, M, \sqrt{-d} \cdot \mathbf{1}]$  generates a self-orthogonal code over  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-d$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.



## Some results...

From permutation representations of  $M_{11}$  on less than 165 points (inclusive), from incidence matrices of weakly self-orthogonal designs we constructed at least 70 non-equivalent non-trivial binary self-orthogonal codes:

- ▶ 6 codes from  $M_{11}$  on 66 points,
- ▶ 14 or more codes from  $M_{11}$  on 110 points,
- ▶ 37 or more codes from  $M_{11}$  on 132 points,
- ▶ 3 or more codes from  $M_{11}$  on 144 points,
- ▶ 10 or more codes from  $M_{11}$  on 165 points.

## Orbit matrices

Let  $\mathcal{D}$  be a  $1 - (v, k, \lambda)$  design and  $G$  be an automorphism group of the design. Let  $v_1 = |\mathcal{V}_1|, \dots, v_n = |\mathcal{V}_n|$  be the sizes of point orbits and  $b_1 = |\mathcal{B}_1|, \dots, b_m = |\mathcal{B}_m|$  be the sizes of block orbits under the action of the group  $G$ . We define an **orbit matrix** as  $m \times n$  matrix:

$$O = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix},$$

where  $a_{ij}$  is the number of points of the orbit  $\mathcal{V}_j$  incident with a block of the orbit  $\mathcal{B}_i$ . It is easy to see that the matrix is well-defined and that  $k = \sum_{j=1}^n a_{ij}$ .

For  $x \in \mathcal{B}_s$ , by counting the incidence pairs  $(P, x')$  such that  $x' \in \mathcal{B}_t$  and  $P$  is incident with the block  $x$ , we obtain

$$\sum_{x' \in \mathcal{B}_t} |x \cap x'| = \sum_{j=1}^m \frac{b_t}{v_j} a_{sj} a_{tj}.$$

Let  $\mathcal{D}$  be a weakly  $q$ -self-orthogonal design such that

$$k \equiv a \pmod{q}$$

and

$$|B_i \cap B_j| \equiv d \pmod{q},$$

for all  $i, j \in \{1, \dots, b\}$ ,  $i \neq j$ , where  $B_i$  and  $B_j$  are two blocks of a design  $\mathcal{D}$ .

Let  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and block orbits of length  $b_1, b_2, \dots, b_m$ , and let  $O$  be an orbit matrix of a design  $\mathcal{D}$  under the action of a group  $G$ .

For  $x \in \mathcal{B}_s$  and  $s \neq t$  it follows that

$$\frac{b_t}{w} O[s] \cdot O[t] \equiv b_t d \pmod{q}, \quad (1)$$

$$\frac{b_s}{w} O[s] \cdot O[s] \equiv a + (b_s - 1)d \pmod{q}. \quad (2)$$

## Codes from orbit matrices of $q$ -self-orthogonal 1-designs

### Theorem ([3])

Let  $\mathcal{D}$  be a self-orthogonal 1-design and  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and block orbits of length  $b_1, b_2, \dots, b_m$  such that  $b_i = 2^o \cdot b'_i$ ,  $w = 2^u \cdot w'$ ,  $o \leq u$ ,  $2 \nmid b'_i, w'$ , for  $i \in \{1, \dots, m\}$ . Then the binary code spanned by the rows of orbit matrix of the design  $\mathcal{D}$  (under the action of the group  $G$ ) is a self-orthogonal code of length  $\frac{v}{w}$ .

### Theorem

Let  $q$  be prime power and  $\mathbb{F}_q$  a finite field of order  $p$ .  
Let  $\mathcal{D}$  be a  $q$  self-orthogonal 1-design and let  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and  $m$  block orbits of length  $w$ . Then the linear code spanned by the rows of orbit matrix of the design  $\mathcal{D}$  (under the action of the group  $G$ ) is a self-orthogonal code over  $\mathbb{F}_q$  of length  $\frac{v}{w}$ .

## Case 2

## Theorem

Let  $\mathcal{D}$  be a weakly self-orthogonal 1-design such that  $k$  is even and the block intersection numbers are odd and  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and block orbits of length  $b_1, b_2, \dots, b_m$  such that  $b_i = 2^o \cdot b'_i$ ,  $w = 2^u \cdot w'$ ,  $o \leq u$ ,  $2 \nmid b'_i, w'$  for  $i \in \{1, \dots, m\}$ . Let  $O$  be the orbit matrix of  $\mathcal{D}$  under action of a group  $G$ .

- a) If  $o = u = 0$ , then the binary linear code spanned by the rows of the matrix  $[I_m, O]$  is a self-orthogonal code of the length  $m + \frac{v}{w}$ .
- b) If  $o \geq 1$  and  $o = u$  then the binary linear code spanned by the rows of the matrix  $[I_m, O, \mathbf{1}]$  is a self-orthogonal code of the length  $m + \frac{v}{w} + 1$ .
- b) If  $o < u$ , then binary linear code spanned by the rows of the matrix  $O$  is a self-orthogonal code of the length  $\frac{v}{w}$ .

Case 2 (over  $\mathbb{F}_q$ )**Theorem**

Let  $q$  be prime power and  $\mathbb{F}_q$  a finite field of order  $p$ .

Let  $\mathcal{D}$  be a weakly  $q$ -self-orthogonal 1-design such that  $k \equiv 0 \pmod{q}$  and  $|B_i \cap B_j| \equiv d \pmod{q}$ , for all  $i, j \in \{1, \dots, b\}$ ,  $i \neq j$ , where  $B_i$  and  $B_j$  are two blocks of a design  $\mathcal{D}$ , and let  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and  $m$  block orbits of length  $w$  and let  $O$  be the orbit matrix of  $\mathcal{D}$  under action of a group  $G$ .

- If  $p \mid w$ , then linear code spanned by the rows of the matrix  $[\sqrt{-d}I_m, O]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $d$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
- If  $p \mid w - 1$ , then linear code spanned by the rows of the matrix  $[\sqrt{wd}I_m, O, \sqrt{-wd}\mathbf{1}]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $wd$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
- If  $p \nmid w$  and  $p \nmid w - 1$ , then linear code spanned by the rows of the matrix  $[\sqrt{wd - (w-1)d}I_m, O, \sqrt{-wd}\mathbf{1}]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-wd$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.

## Case 3

**Theorem ([3])**

Let  $\mathcal{D}$  be a weakly self-orthogonal 1-design such that  $k$  is odd and the block intersection numbers are even and  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and block orbits  $b_1, b_2, \dots, b_m$  such that  $b_i = 2^o \cdot b'_i$ ,  $w = 2^u \cdot w'$ ,  $o \leq u$ ,  $2 \nmid b'_i, w'$ , for  $i \in \{1, \dots, m\}$ . Let  $O$  be the orbit matrix of  $\mathcal{D}$  under action of a group  $G$ .

- If  $o = u$ , then the binary linear code spanned by the rows of matrix  $[I_m, O]$  is a self-orthogonal code of length  $m + \frac{v}{w}$ .
- If  $o < u$ , then the binary linear code spanned by the rows of matrix  $O$  is a self-orthogonal code of length  $\frac{v}{w}$ .

**Theorem**

Let  $q$  be prime power and  $\mathbb{F}_q$  a finite field of order  $p$ .

Let  $\mathcal{D}$  be a weakly  $q$ -self-orthogonal design such that  $k \equiv a \pmod{q}$  and block intersection numbers are multiples of  $q$ , and let  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and  $m$  block orbits of length  $w$ . Then the linear code spanned by the rows of matrix  $[\sqrt{-a}I_m, O]$ , where  $O$  is orbit matrix of the design  $\mathcal{D}$  (under the action of the group  $G$ ), is a self-orthogonal code over  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $a$  is a square in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.

## Case 4

**Theorem**

Let  $\mathcal{D}$  be a weakly self-orthogonal 1-design such that  $k$  is odd and the block intersection numbers are odd and  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and block orbits of length  $b_1, b_2, \dots, b_m$  such that  $b_i = 2^o \cdot b'_i$ ,  $w = 2^u \cdot w'$ ,  $o \leq u$ ,  $2 \nmid b'_i, w'$ , for  $i \in \{1, \dots, m\}$ . Let  $O$  be the orbit matrix of  $\mathcal{D}$  under action of a group  $G$ .

- If  $o = u = 0$ , then the binary linear code spanned by the rows of the matrix  $[O, \mathbf{1}]$  is a self-orthogonal code of the length  $\frac{v}{w} + 1$ .
- Otherwise, the binary linear code spanned by the rows of the matrix  $O$  is a self-orthogonal code of the length  $\frac{v}{w}$ .



## Theorem

Let  $q$  be prime power and  $\mathbb{F}_q$  a finite field of order  $q$ . Let  $\mathcal{D}$  be a  $1 - (v, k, r)$  design such that  $k \equiv a \pmod{q}$  and  $|B_i \cap B_j| \equiv d \pmod{q}$ , for all  $i, j \in \{1, \dots, b\}$ ,  $i \neq j$ , where  $B_i$  and  $B_j$  are two blocks of a design  $\mathcal{D}$ , and let  $G$  be an automorphism group of the design which acts on  $\mathcal{D}$  with  $n$  point orbits of length  $w$  and  $m$  block orbits of length  $w$  and let  $O$  be the orbit matrix of  $\mathcal{D}$  under action of a group  $G$ .

- ▶ If  $a = d$  we differ two cases.
  - a) If  $p \mid w$ , then linear code spanned by the rows of the matrix  $O$  is a self-orthogonal code over the field  $\mathbb{F}_q$ .
  - b) If  $p \nmid w$ , then linear code spanned by the rows of the matrix  $[\sqrt{-a}I_m, O]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-a$  is square root in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
- ▶ If  $a \neq d$ , we differ three cases.
  - a) If  $p \mid w$ , then linear code spanned by the rows of the matrix  $[\sqrt{d-a}I_m, O]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $d-a$  is square root in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
  - b) If  $p \mid w - 1$ , then linear code spanned by the rows of the matrix  $[\sqrt{wd-a}I_m, O, \sqrt{-wd}\mathbf{1}]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-wd$  is square root in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.
  - c) If  $p \nmid w$  and  $p \nmid w - 1$ , then binary linear code spanned by the rows of the matrix  $[\sqrt{d-a}I_m, O, \sqrt{-wd}\mathbf{1}]$  is a self-orthogonal code over the field  $\mathbb{F}$ , where  $\mathbb{F} = \mathbb{F}_q$  if  $-wd$  is square root in  $\mathbb{F}_q$ , and  $\mathbb{F} = \mathbb{F}_{q^2}$  otherwise.

## Some results...

From permutation representations of  $M_{11}$  on less than 165 points (inclusive), from orbit matrices we constructed at least 87 non-equivalent non-trivial binary self-orthogonal codes:

- ▶ 2 codes from  $M_{11}$  on 66 points,
- ▶ 22 codes from  $M_{11}$  on 110 points,
- ▶ 21 codes from  $M_{11}$  on 132 points,
- ▶ 24 or more codes from  $M_{11}$  on 144 points,
- ▶ 18 or more codes from  $M_{11}$  on 165 points.

8 of constructed codes are optimal:

$[10, 4, 4]$ ,  $[12, 5, 4](2)$ ,  $[12, 6, 4]$ ,  $[12, 11, 2]$ ,  $[16, 5, 8]$ ,  $[24, 12, 8]$ ,  $[31, 15, 8]$  and one of them is best known:  $[96, 48, 16]$ .

Thank you for your attention!